Lecture 11 *Linear programming : The Revised Simplex Method*

11.1 The Revised Simplex Method

While solving linear programming problem on a digital computer by regular simplex method, it requires storing the entire simplex table in the memory of the computer table, which may not be feasible for very large problem. But it is necessary to calculate each table during each iteration. The revised simplex method which is a modification of the original method is more economical on the computer, as it computes and stores only the relevant information needed currently for testing and / or improving the current solution. i.e. it needs only

- The net evaluation row Δ_i to determine the non-basic variable that enters the basis.
- The pivot column
- The current basis variables and their values $(X_B \text{ column})$ to determine the minimum positive ratio and then identify the basis variable to leave the basis.

The above information is directly obtained from the original equations by making use of the inverse of the current basis matrix at any iteration.

There are two standard forms for revised simplex method

- **Standard form-I** In this form, it is assumed that an identity matrix is obtained after introducing slack variables only.
- **Standard form-II** If artificial variables are needed for an identity matrix, then twophase method of ordinary simplex method is used in a slightly different way to handle artificial variables.

11.2 Steps for solving Revised Simplex Method in Standard Form-I

```
Solve by Revised simplex method
Max Z = 2x_1 + x_2Subject to
       3x_1 + 4x_2 \le 66 x_1 + x_2 \leq 3and x_1, x_2 \geq 0
```
SLPP

```
Max Z = 2x_1 + x_2 + 0s_1 + 0s_2Subject to
        3 x_1 + 4 x_2 + s_1 = 66 x_1 + x_2 + s_2 = 3and x_1, x_2, s_1, s_2 \ge 0
```
Step 1 – Express the given problem in standard form -1

- Ensure all $b_i \geq 0$
- The objective function should be of maximization

 Use of non-negative slack variables to convert inequalities to equations The objective function is also treated as first constraint equation

 $Z - 2x_1 - x_2 + 0s_1 + 0s_2 = 0$ $3 x_1 + 4 x_2 + s_1 + 0s_2 = 6$ -- (1) $6 x_1 + x_2 + 0s_1 + s_2 = 3$ and $x_1, x_2, s_1, s_2 \ge 0$

Step 2 – Construct the starting table in the revised simplex form Express (1) in the matrix form with suitable notation

Column vector corresponding to Z is usually denoted by e_1 matrix B₁, which is usually denoted as B₁ = $[\beta_0^{(1)}, \beta_1^{(1)}, \beta_2^{(1)} \dots \beta_n^{(1)}]$

Hence the column $\beta_0^{(1)}$, $\beta_1^{(1)}$, $\beta_2^{(1)}$ constitutes the basis matrix B_1 (whose inverse B_1^{-1} is also B_1)

Step 3 – Computation of Δ_j for $a_1^{(1)}$ and $a_2^{(1)}$

 Δ_1 = first row of B₁⁻¹ * a₁⁽¹⁾ = 1 * -2 + 0 * 3 + 0 *6 = -2 Δ_2 = first row of B₁⁻¹ * a₂⁽¹⁾ = 1 * -1 + 0 * 4 + 0 *1 = -1

Step 4 – Apply the test of optimality

Both Δ_1 and Δ_2 are negative. So find the most negative value and determine the incoming vector.

Therefore most negative value is $\Delta_1 = -2$. This indicates $a_1^{(1)}(x_1)$ is incoming vector.

Step 5 – Compute the column vector X_k

$$
X_k = B_1^{-1} * a_1^{(1)}
$$

$$
\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \ast \begin{bmatrix} -2 \\ 3 \\ 6 \end{bmatrix} = \begin{bmatrix} -2 \\ 3 \\ 6 \end{bmatrix}
$$

Step 6 – Determine the outgoing vector. We are not supposed to calculate for Z row.

	B_1					
Basic variables	e ₁ (Z)	$\beta_1^{(1)}$	$\beta_2^{(1)}$	X_B	X_k	X_B/X_k
7					-7	
S_1				n		
S_2					6	$1/2 \rightarrow$ outgoing
					incoming	

Step 7 – Determination of improved solution

Column e_1 will never change, x_1 is incoming so place it outside the rectangular boundary

Make the pivot element as 1 and the respective column elements to zero.

Construct the table to start with second iteration

 $\Delta_4 = 1 * 0 + 0 * 0 + 1/3 * 1 = 1/3$ $\Delta_2 = 1 * -1 + 0 * 4 + 1/3 * 1 = -2/3$ Δ_2 is most negative. Therefore a_2 ⁽¹⁾ is incoming vector.

Compute the column vector

Determine the outgoing vector

Determination of improved solution

 $\Delta_4 = 1 * 0 + 4/21 * 0 + 5/21 * 1 = 5/21$ $\Delta_3 = 1 * 0 + 4/21 * 1 + 5/21 * 0 = 4/21$

 Δ_4 and Δ_3 are positive. Therefore optimal solution is Max Z = 13/7, x₁= 2/7, x₂ = 9/7

11.3 Worked Examples

Example 1

```
Max Z = x_1 + 2x_2Subject to
        x_1 + x_2 \leq 3x_1 + 2x_2 \leq 53x_1 + x_2 \le 6and x_1, x_2 \ge 0
```
Solution

SLPP

 $Max Z = x_1 + 2x_2 + 0s_1 + 0s_2 + 0s_3$ Subject to $x_1 + x_2 + s_1 = 3$ $x_1 + 2x_2 + s_2 = 5$ $3x_1 + x_2 + s_3 = 6$ and $x_1, x_2, s_1, s_2, s_3 \ge 0$

Standard Form-I

 $Z - x_1 - 2x_2 - 0s_1 - 0s_2 - 0s_3 = 0$ $x_1 + x_2 + s_1 + 0s_2 + 0s_3 = 3$ $x_1 + 2x_2 + 0s_1 + s_2 + 0s_3 = 5$ $3x_1 + x_2 + 0s_1 + 0s_2 + s_3 = 6$ and $x_1, x_2, s_1, s_2, s_3 \ge 0$

Matrix form

Revised simplex table Additional table

Computation of Δ_j for $a_1^{(1)}$ and $a_2^{(1)}$

 Δ_1 = first row of B₁⁻¹ * a₁⁽¹⁾ = 1 * -1 + 0 * 1 + 0 *1 + 0 *3= -1 Δ_2 = first row of B₁⁻¹ * a₂⁽¹⁾ = 1 * -2 + 0 * 1 + 0 *2+ 0 *1 = -2

 Δ_2 = -2 is most negative. So a_2 ⁽¹⁾(x₂) is incoming vector.

Compute the column vector X_k

Improved Solution

Revised simplex table for II iteration

 $\Delta_1 = 1 * -1 + 0 * 1 + 1 *1 + 0 *3 = 0$ $\Delta_4 = 1 * 0 + 0 * 0 + 1 * 1 + 0 * 0 = 1$

 Δ_1 and Δ_4 are positive. Therefore optimal solution is Max Z = 5, x₁= 0, x₂ = 5/2